

## Background

In this work, we create circuits with extra ancillas to protect the logical exponential operator by a quantum error detecting code.

- **Assumptions.** We consider Pauli noise after gates and hardware with full connectivity and mid-circuit measurements.
- **Notation.** We use  $\overline{\mathcal{P}}$  for the logical operator and  $\mathcal{P}$  to represent the physical operator corresponding to  $\overline{\mathcal{P}}$ .  $I, X, Y, Z$  are single qubit Pauli operators.

## Product Formula Simulation

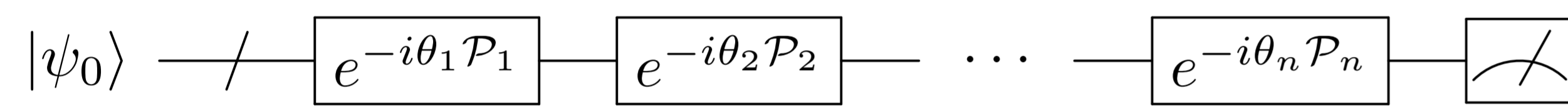
- For fermionic system simulation, convert the Hamiltonian into Pauli form using Jordan-Wigner, Bravyi-Kitaev or parity mapping

$$H \xrightarrow{\text{mapping}} \tilde{H}, \text{ where } \tilde{H} = \sum_j \alpha_j \mathcal{P}_j \quad (1)$$

- Use product formula approach (Trotter-Suzuki decomposition) to decompose

$$e^{-i\theta\tilde{H}} \simeq \prod_j e^{-i\theta_j\mathcal{P}_j}, \text{ where } \theta_j = \alpha_j\theta \quad (2)$$

- Create circuits in the following form:



## Quantum Error Detection Code

The stabilizer generators are  $X^{\otimes n}$  and  $Z^{\otimes n}$ . Logical Pauli operators are given by  $\overline{X}_i = X_i X_{n-1}, \overline{Z}_i = Z_i Z_n$  and  $\overline{Y}_i = i\overline{X}_i \overline{Z}_i$ . The logical rotation gates are

$$\overline{R}_{X_i}(\theta) = \exp\left(-i\frac{\theta}{2} X_i X_{n-1}\right), \overline{R}_{Z_i}(\theta) = \exp\left(-i\frac{\theta}{2} Z_i Z_{n-1}\right) \quad (3)$$

Below are examples of logical rotations under the  $[[4, 2, 2]]$  code.

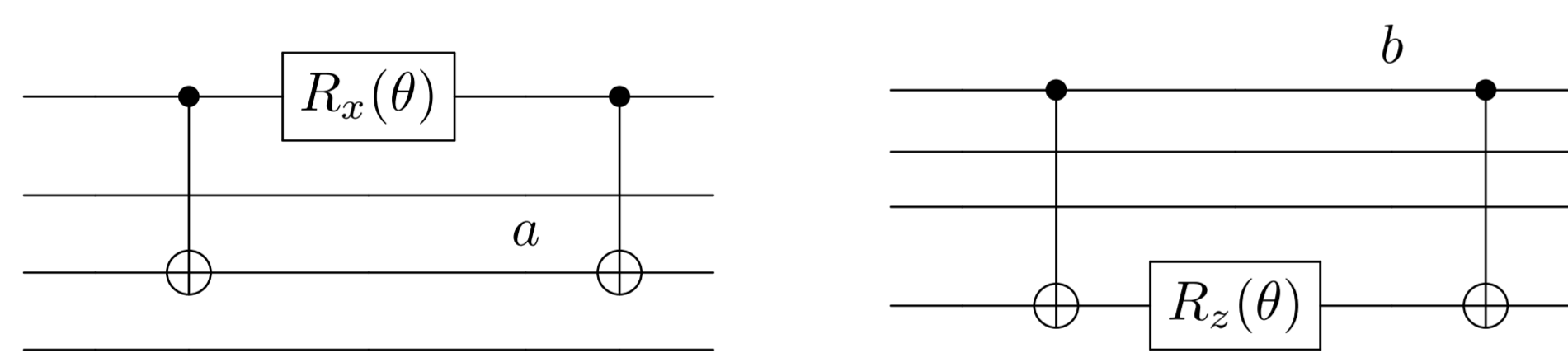


Figure 1. Circuit of  $\overline{R}_{X_1}(\theta) = \exp(-i\theta X_1 X_3/2)$  (left) and  $\overline{R}_{Z_1}(\theta) = \exp(-i\theta Z_1 Z_4/2)$  (right).

They are not fault-tolerant, since a  $Z_3$  error occurring at location  $a$  becomes a  $Z_2 Z_4$  error, and a  $X_1$  error at location  $b$  becomes a  $X_1 X_4$  error, both undetectable.

## Weakly Fault Tolerant Rotation

Previous research (not yet published) from Todd Brun and Christopher Gerhard purposes a weakly fault-tolerant construction shown in Fig. 2 that helps detect most single- and two-qubit Pauli noise, except for  $X$  error at location  $d$  and  $Z$  error at location  $f$ .

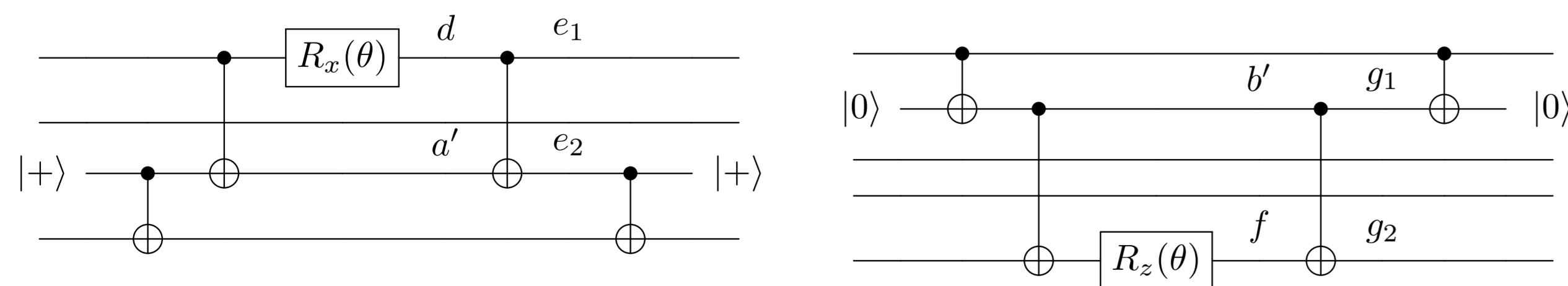


Figure 2. Circuit of weakly fault tolerant  $\overline{R}_{X_1}(\theta)$  (left) and  $\overline{R}_{Z_1}(\theta)$  (right)

## Approach

Four steps to protect each exponential map  $\exp(-i\theta\overline{\mathcal{P}})$  with  $[[n, n-2, 2]]$  code.

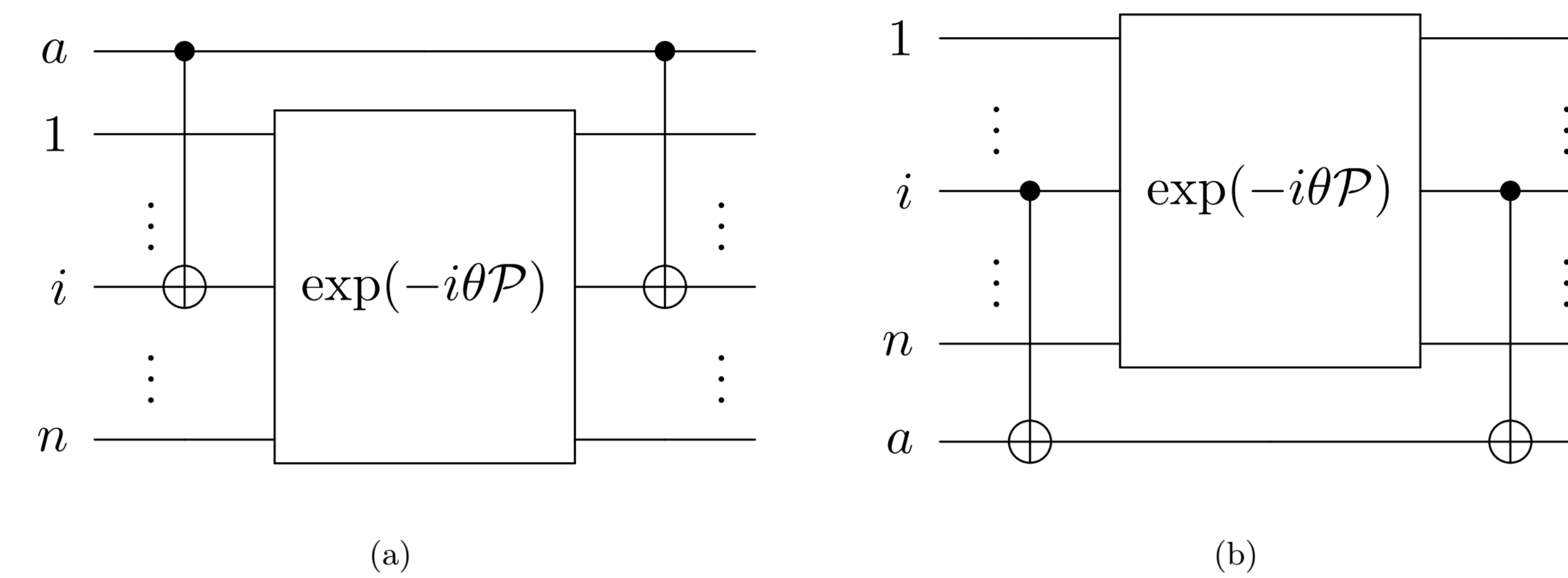
### Step 1: Construct Physical Operator

Construct physical operator by  $\exp(-i\theta\overline{\mathcal{P}}) \rightarrow \exp(-i\theta\mathcal{P})$ , which fulfill two requirements.

- The given circuit should perform logical operation  $\overline{U}$  on logical state  $|\overline{\psi}\rangle$
- The given circuit should always convert any valid logical state (i.e., a state that falls into the codespace) into another valid logical state in same subspace.

### Step 2: Create Equivalent Circuits

Create a group of circuits that equivalently implement  $\exp(-i\theta\mathcal{P})$  with one (or more) ancilla. Use the following construction to create equivalent circuits:



Here the  $i$ -th qubit **can be any qubits** between qubit 1 and qubit  $n$ . The new map  $U$  is still an exponential map  $\exp(-i\theta\mathcal{P}')$ , and we can set every possible qubit as ancilla.

### Step 3: Search for Noise-Resilient Circuit

The most-noise resilient circuit is the one with the least logical error ratio ( $\text{num\_logical\_error}/\text{num\_total\_err}$ ) under all possible depolarizing Pauli noise after each gate. Use numerical simulation to find the noise-resilient circuit.

### Step 4: Perform Mid-circuit syndrome Measurement

Use noise-resilient operator construction for  $\exp(-i\theta\mathcal{P})$ . After measuring error syndrome after each  $\exp(-i\theta\mathcal{P})$ , we discard the noisy shots and use the remaining shots for further analysis.

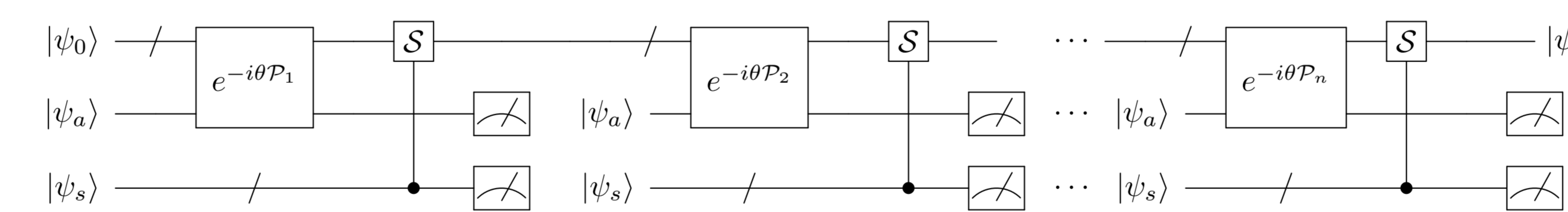


Figure 3. Mid-Circuit Measurement Protocol.

## Comments

1. Advantages of our approach:
  - Easy to implement in near-term devices.
  - Optimize near term resources by using fewer qubits and  $\mathcal{O}(2n)$  CNOTs.
  - Reduce overhead by discarding noisy shots and no need for correction compare with QEC.
  - Can be used on near-term small-scale fermionic system simulation.
  - It can be easily extended for any  $[[n, k, d]]$  stabilizer code.
2. Disadvantages of using our approach
  - Required device with full connectivity to achieve optimal performance (e.g. ion trap system)
  - Required device with perfect qubit reset to achieve best mid-circuit measurement.

## Examples and Performance

### Logical Rotation-Y

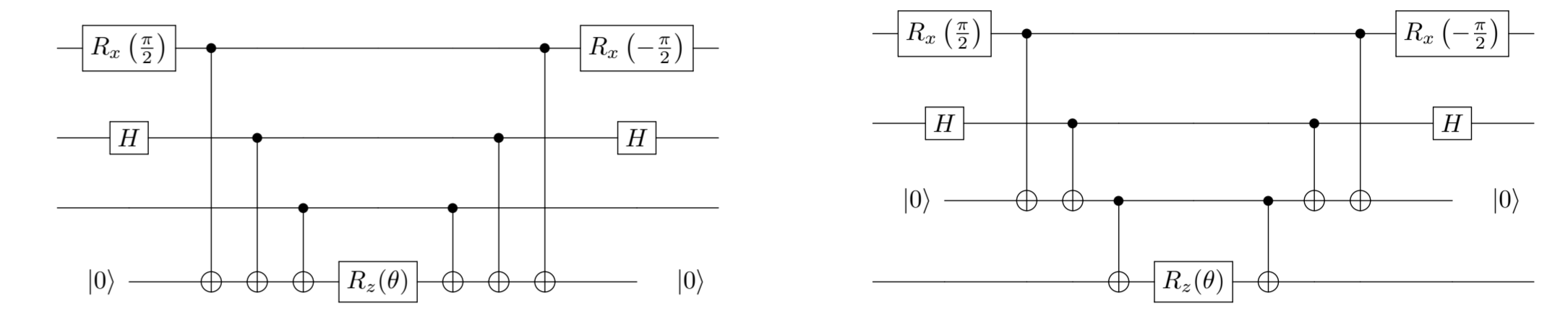


Figure 4. Two noise-resilient circuit that implements  $\overline{R}_{Y_i}(\theta) = \exp(-i\theta Y_i X_{n-1} Z_n/2)$ .

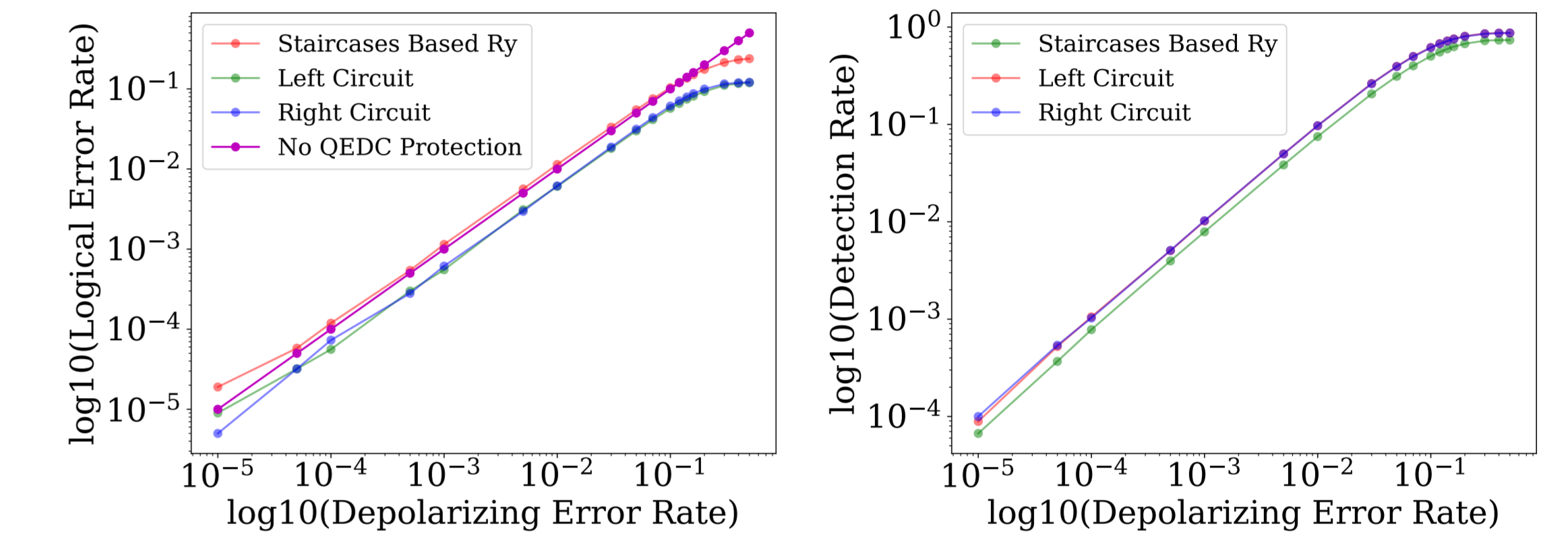


Figure 5. Logical error rate (left) and post-selection rate (right) of two noise-resilient circuits under 1,000,000 trials simulation with depolarizing gate noise.

### Logical Exponential with Weight $n > 3$ Pauli

Consider noise-resilient circuits for  $\exp(-i\theta Z^{\otimes n})$  with different weights  $n$ .

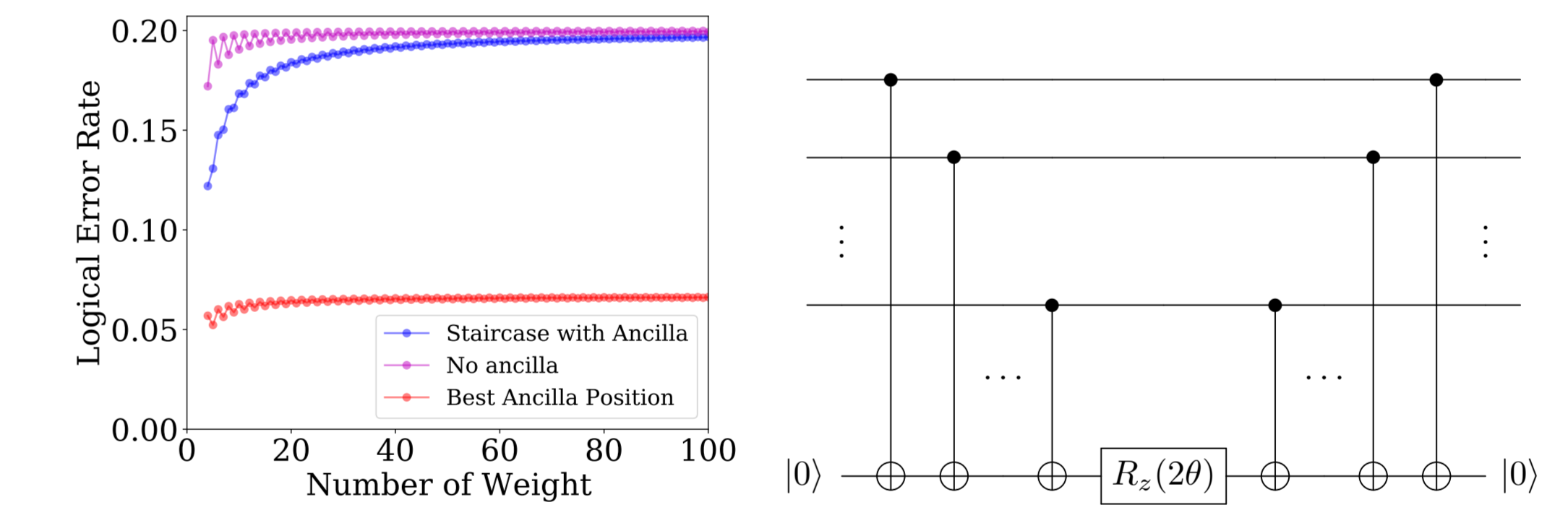


Figure 6. Logical error rate versus weights in  $\exp(-i\theta Z^{\otimes n})$  using circuit on the right.

### Performance in Deep Circuit

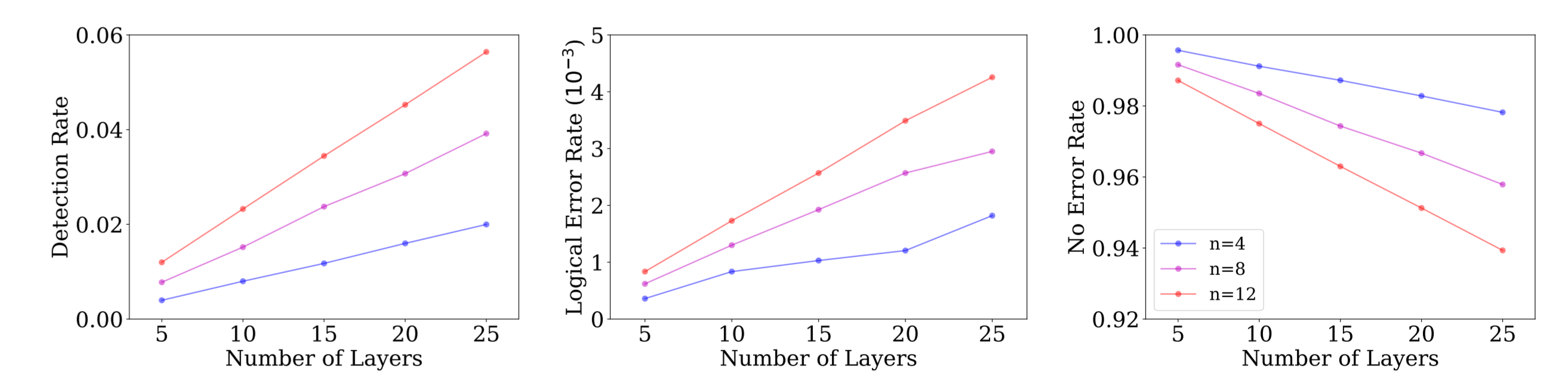


Figure 7. Detection rate, logical error rate and no error rate with different number of  $\exp(-i\theta Z^{\otimes n})$  layers and qubit  $n$  under error probability  $10^{-4}$ .

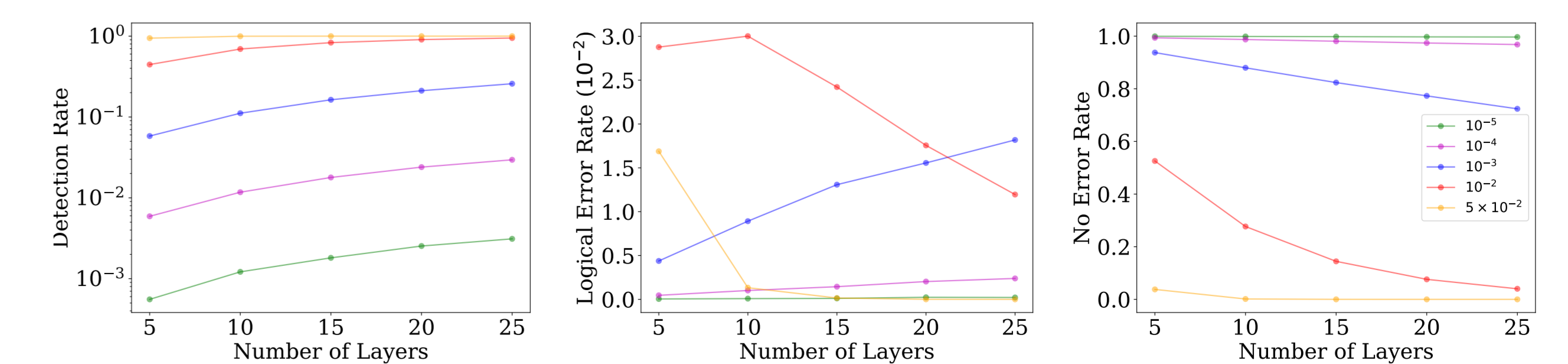


Figure 8. Detection rate, logical error rate and no error rate with different number of  $\exp(-i\theta Z^{\otimes 6})$  layers and error probability.